

DBMS - 3

1 - Database funda.

2 - ERD + Data integ.

3 - Data Normalisation

- Anomalies
- Functional Deps
- Normal Forms

Transaction

- Intro
- ACID properties

<u>ID</u>	NAME	BATCH_ID	BATCH_NAME
		NULL	

1	John	1	Shenlock
2	Mary	1	Shenlock
3	Kilvish	2	Shakti
4	NULL	3	Rebellion
5	Tontia	NULL	NULL

How can I add a student w/o a batch X

- ① Cannot create batch w/o student
- ② Cannot create student w/o batch

In section [anomalies]

issues

- ① missing data (batch w/o student)
- ② FK constraints (student w/o batch)

Deletion anomaly

ID	Name	Batch
1	John	Shenlock
2	M...	...

	2	Primary	Shorlock
→	3	Kilvish	Shobhi

Shobhi

None | PSD | Agency

Deletion anomaly

↳ bad design

Update Anomaly

	ID	Name	BID	BName
→	1	John	1	<u>Shorlock</u>

→ 2	Mong	1	Shenlock
→ 3	Mycroft	1	Shenlock

Update Shenlock to Shenlock
Season 6

Failure

- manual
- automated

Inconsistent data

Common issues in DB

→ bad design

Insertion

↳ missing data

↳ FK constraints

Deletion

↳ Tightly couple

Update

↳ errors or partial failures

Good design - Normalisation

Functional dependencies

CAR



VIN → Engine Capacity

VIN → Model

Engine Capacity → VIN



ID	NAME	PHONE	BATCH_ID	B_NAME
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(ID) → Name

(ID) → Phone

ID → ... }

BATCH_ID → BATCH_NAME

NAME → (ID)



NOT Unique

Phone → ID

ID → Phone

Select phone from
STU where ID=4

ID → Batch ID
↙ ↘

ratings, feedback

Session Queries

Student ID	Main ID	Session ID	Rating	Feedback
①	①	1		

①	①	2	
2	①	3	
2	1	4	

PK - Student, ^{PK} Mentor, SessionID

(FD) 1 Student can only have 1 mentor

Student ID → Mentor ID

2 { StudentID, MentorID, SID }

→ noting
→ feedback

Normalisation

Multiple tables

↳ reduce data redundancy

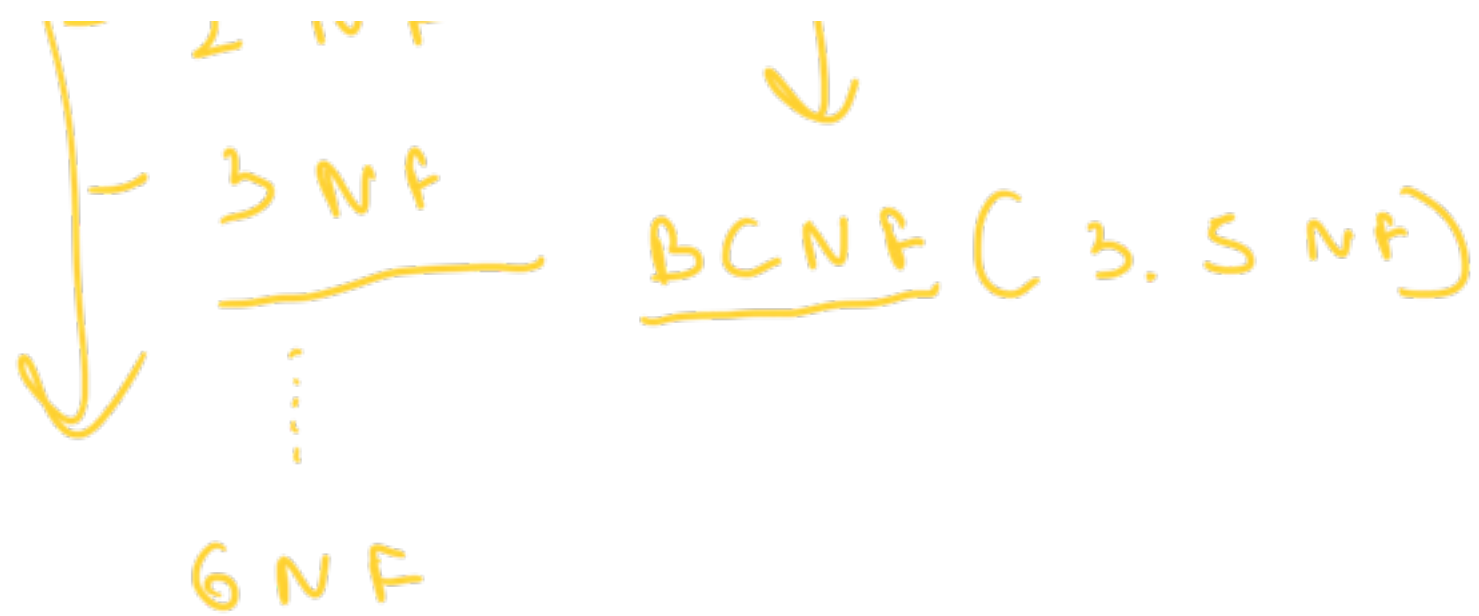
↳ remove anomalies

→ reduce redundancy

→ improve data integrity

Normal Forms (Relational)

	Strictness	DB
1NF	1	
2NF	1	



Denormalisation

2NF to 3NF

Set of rules to normalise your database

1NF

2NF - 1NF + new rules

1NF

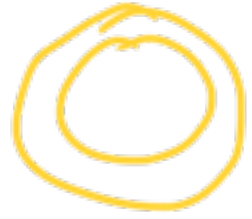
Rule

→ All your values should be

atomic

ER model

- multivalued



- composite

id:

} collections

Mysql 7.5.7

1NF

- X collections

ID	NAME	PHONE_NOS	X 1NF
----	------	-----------	-------

1	Sherlock	[12345, ...]	}	Lists
2	Morant	[98765, ...]		

① Multiple columns

ID	NAME	Phone #1	Phone #2	...
1	Sherlock	1234	9876 ...	

INF ✓
 .m a -

③

① We don't know the upper bound

② Efficient

↳ Find a student by phone

→ All phone no.

③ Sparse

X

② Multiple rows

INF ✓

ID	Name	phone
1	Shenlock	12345
1	Shenlock	98765..
2	Soh	123456..

} X

2 | John | 9876543210

① redundant data (n-1)

2 b a

② PK - ID + Phone

Phone

CSV a, b, c

"12345; 98765;"

"123, 987, ..."

var char

□

Search user when

no = (123)

Complicated solution

join on phone
deleting

Occom's Roger



Simplest solution

KZSS

Simplest - Create a new table

New Table

SID	Name	Email	<u>Student</u>
-----	------	-------	----------------

Phone

S ID	Phone number
1	12 345
1	987654 ..

Normalise

User

Normal form

ID	NAME	{interests}
1	Tony	{DBMS } *20 30



→ Multiple columns ✓

✓ Multiple Rows ✓

→ Multiple rows

ID Phone No

1 123

1 987

→ Create a new table

Students

Phone numbers



MC

User ID

Name

Interest 1

Interest 2



Separate table

Student ID	Interest
1	DBMS
1	Books
1	Memes

User interest

↳ 1 table (user)

↳ 4 tables

→ performance worse

→ optimisations

6:12

10:42

6:17

10:47

4 sessions

6 sessions

4 sessions

— Theory

5th → Queries

5th → Sequences

6th → Transactions

→ Joins

→ Views

→ Triggers

→ Procedures

— ...

2:30

Transactions

→ window

→ Query Optimization



2NP

fanmay. Racker

SOL Fk & dte

2NF

① 1NF

② No partial dependencies

Ⓐ B

A → X

· ID | NAME | BATCH_ID | B_NAME | PSP

ID → NAME

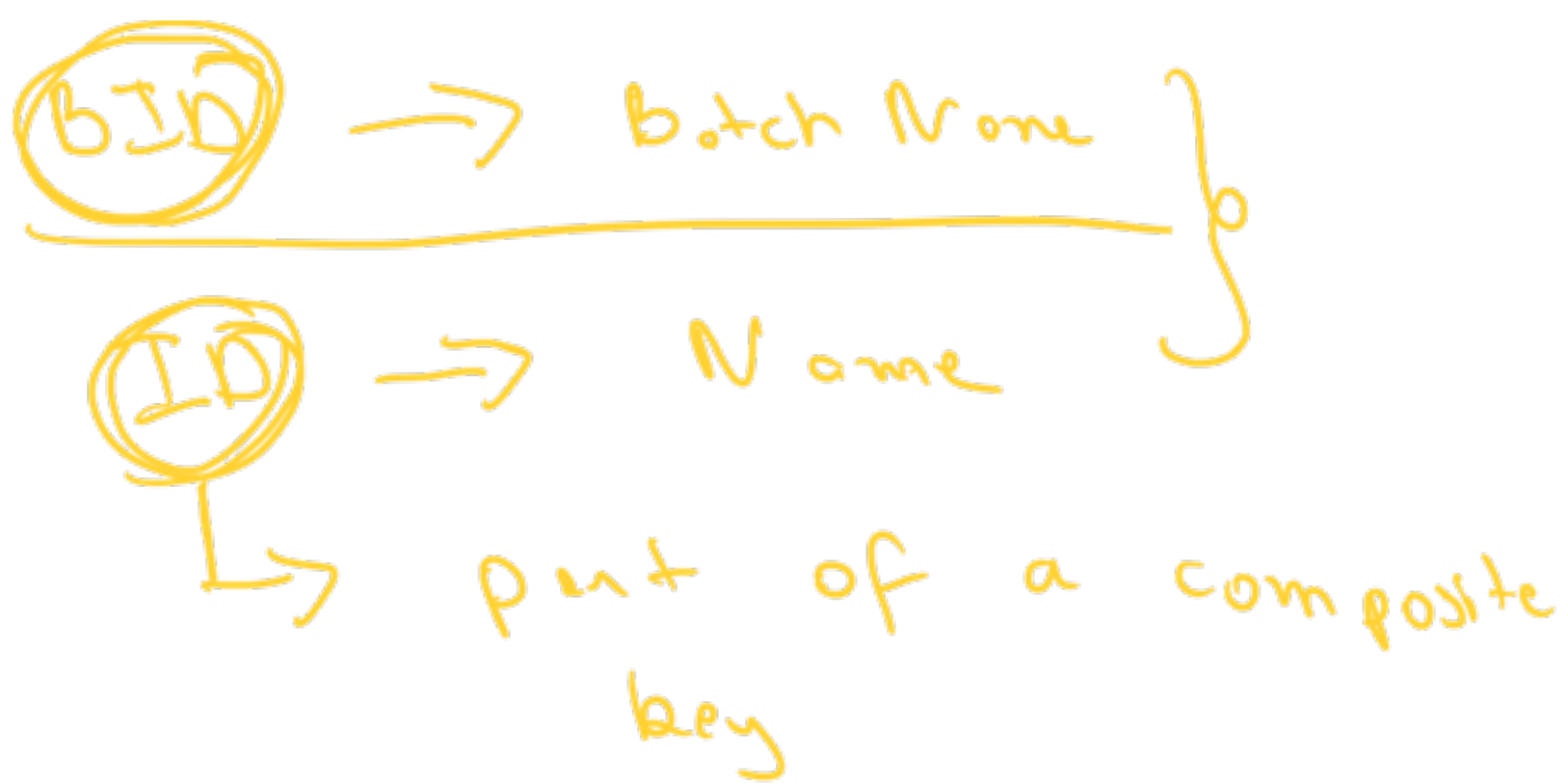
One user can be part of multiple batches

ID	Name	bID	bName	PSP
1	Shenlock	1	Season 1	90%
1	Shenlock	2	Season 2	100%

PK — ID + bID

{ID, bID} → PSP

Functional dep.



Parts of CPK on LHS
partial dependency



2NF

→ Not have partial FD

→ Every non-candidate key attribute should depend on the whole candidate key
* just not a part of it

2NF - Composite

1NF + single PK = 2NF

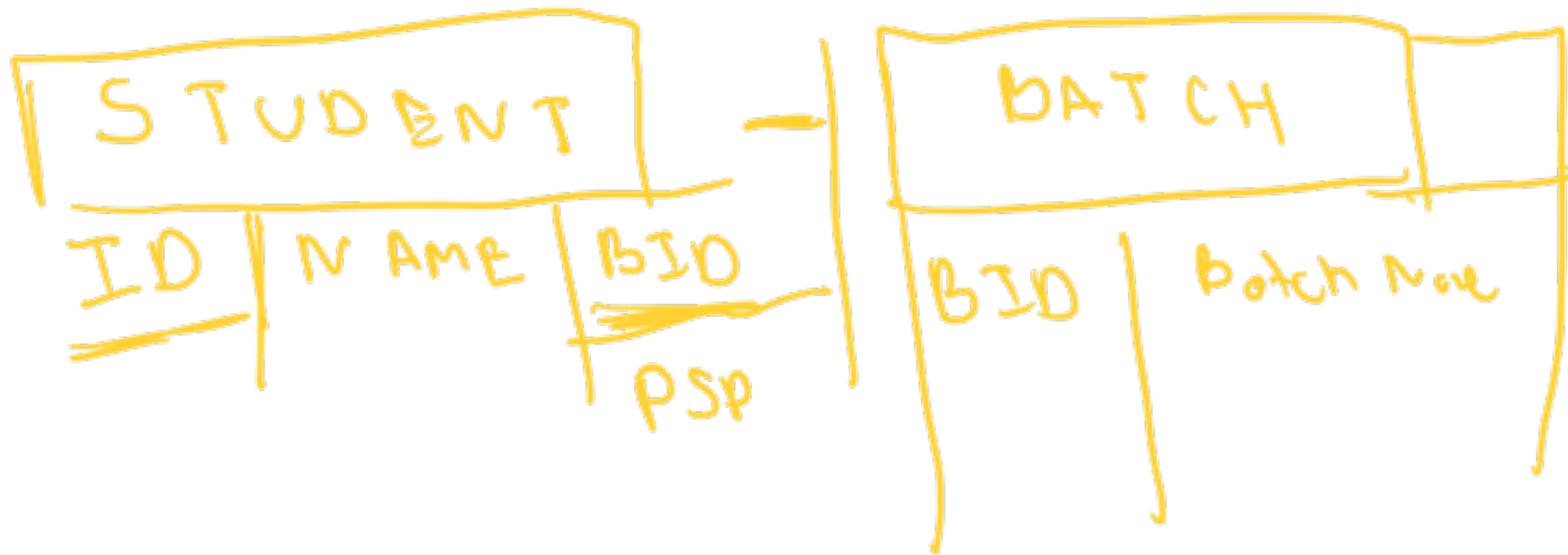
→ composite = \overline{PFD}

ID	NAME	BID	BNAME	PSP
1	Tony	1	X	Y



BID → BNAME - PFD

PK ⇒ ID, BID



2NF - Composite key

Partial Functional Dependency





How to fix)

Create a new table with duplicate involves

① Student can be part of multiple batches

IDEAL

STUDENT	<u>ID</u>	<u>BID</u>	PSP
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1	2	90
---	---	----

1	1	100
---	---	-----

BID → BNAME

ID → NAME

↳ User info

3NF

- 2NF

- It should not contain
transitive dependencies

Transitivity

$a = b$

$b = c$

$a = c$

$X \rightarrow Y$ $Y \rightarrow Z$
 $X \rightarrow Z$ ✓

<u>ID</u>	BID	BNAME	NAME
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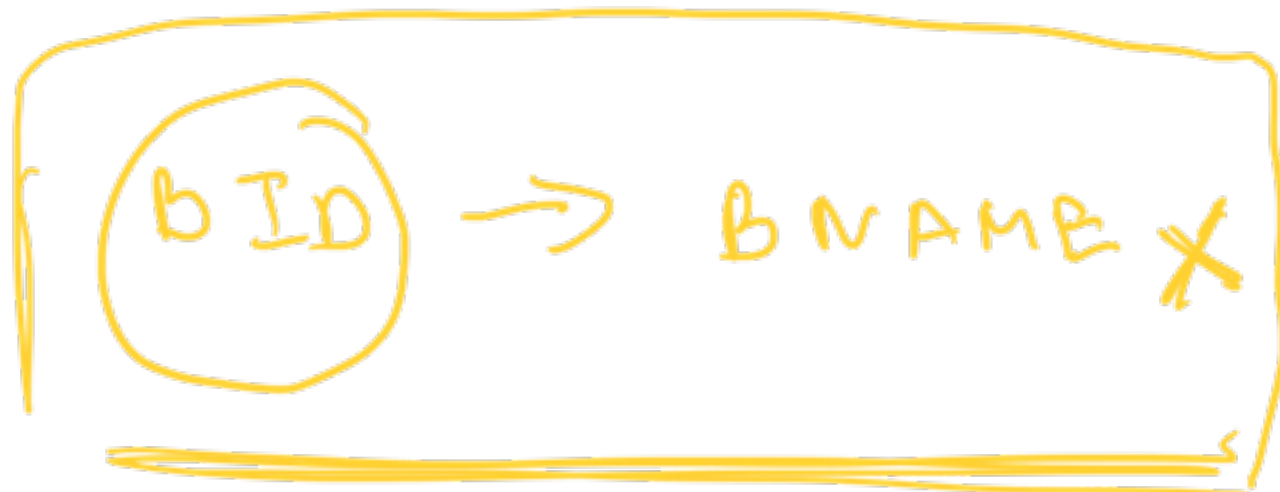


✗ 3NF



Should be in a
candidate

ID	NAME	PHONE	BID	BNAME
----	------	-------	-----	-------



3NF

① No transitive dependencies

ID → BID → BNAME ✗

②

X → Y → ^{(A)(B)}
primary attribute

↓
Super key

3NF → No transitive dependencies

ID → BID → BNAME ✗

A → B → C

A → C

→ ...



BCNF

Boyce Codd NF

3.5 NF

relational model

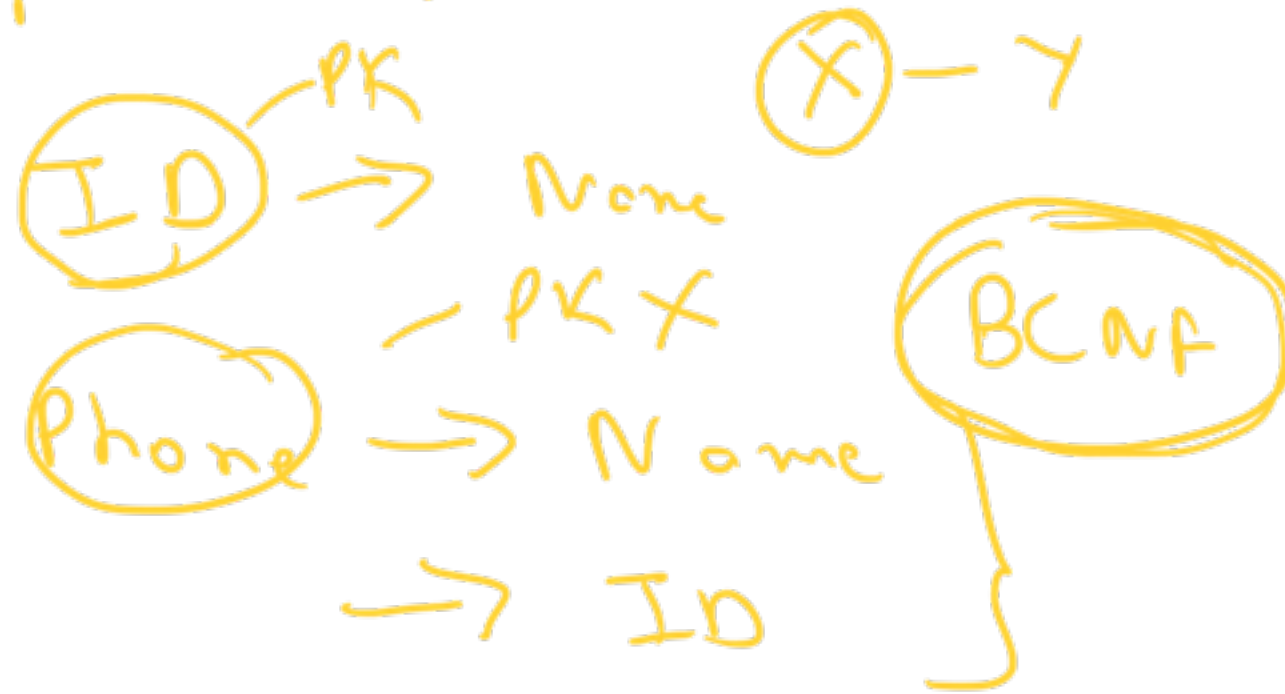


X BCNF

FD

X is not PK

ID	Phone	Name	Batch ID
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Phone — SK, CK, PK (ID)

BCNF → 3NF
 n NF → (n-1) NF
 → nth rule to other PK in X

ID	NAME	Phone	BID	BNAME
----	------	-------	-----	-------

Students

ID	NAME	(BID)	STID	Phone No.
1	Tony		1	123
			1	987



BID	BNAME
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Normalization

→ Check if rules are met

→ If not, create the rule

• Data anomalies

→ Insertion

→ Deletion

→ Update

loss of data

→ errors

inconsistent

• Functional Dependency



ID → NAME



→ Partial Popularity

Composite keys $(\bar{A})B$

$A \rightarrow X$
 $B \rightarrow Y$

→ Transitive

$a = b = c$

$a = c$

$X \rightarrow Y \quad Y \rightarrow Z$

$X \rightarrow Z$

$id \rightarrow BID \rightarrow BNAME$

1NF - Atomic

2NF - No PRD

3NF - No TFD

BCNF = Only PK on LHS of FD

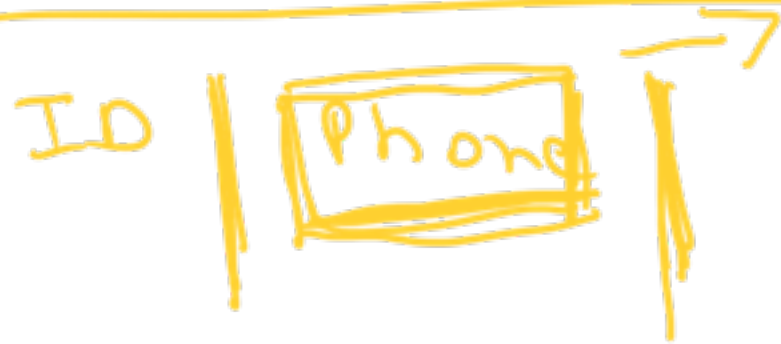


Create a separate table

- Transactions + ACID
- Indexes

NF - (redundancy)

inconsistency



3NF \rightarrow Transitive FD

① X should be a super key

② Y should be a prime attribute

Y = (A)(B)

X(A) Z B



{ StudentID }

{ phone }

{ StudentID, phone }



prime attribute

3NF $X \rightarrow SK$

$Y \rightarrow PA$

2 Student ID, Phone

present in a candidate key

2 SID, (Name) $\rightarrow PA$

SID \rightarrow PK \rightarrow CK \checkmark

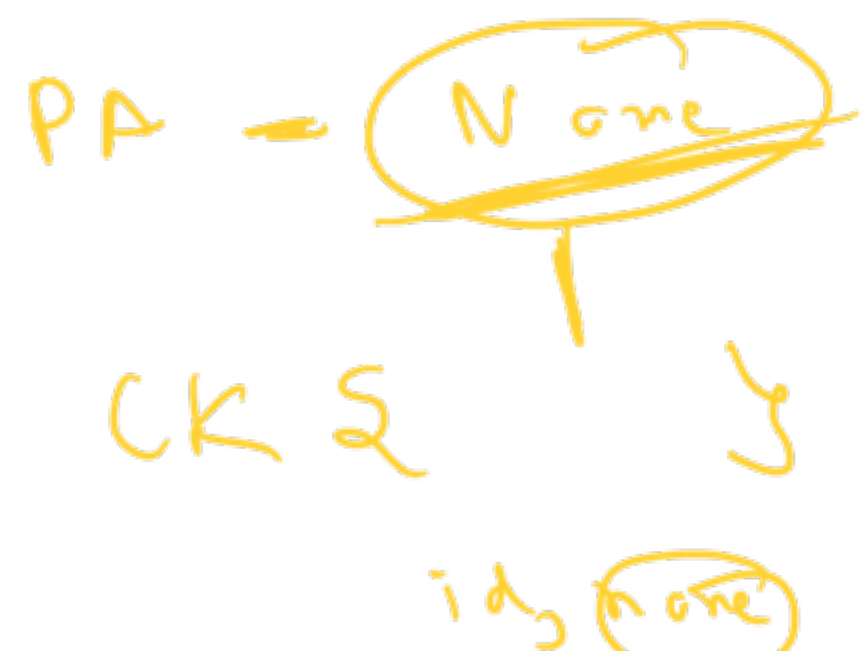
\rightarrow (CK)

S (SID) (Phone)



3NF
 - X - SK }
 - Y - PA }

SID - PK
 - PA



CK - $\{ \textcircled{A}, \textcircled{B} \} \rightarrow X$

$\{ A, B, C \}$

$\{ A, B \} \rightarrow A$

3NF

\rightarrow TFD

BCNF - 3SNF

$\textcircled{A} \rightarrow X$
 $\textcircled{A} \rightarrow \textcircled{PK}$

ID | NAME | Phone

PK ✓
ID → None
ID → Phone

Phone → ID

Phone → None

X BCNF

ID | NAME

ID | PHONE